

# Give Me Letters 2, 3 and 6!

## Partial Password Implementations and Attacks

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# Outline

Partial Passwords

Survey

Guessing Attacks

Recording Attacks

Summary

# Partial Passwords

## Definitions and examples

A **partial password** is a challenge on a subset of characters from a full password.

A **partial password scheme** is an authentication system using partial passwords.

Card Number: XXXX-XXXX-XXXX-XXXX

Personal Greeting: Welcome to SecureCode

Login: XXXX-XXXX-XXXX

Enter the fourth, fifth and sixth characters of your SecureCode:

[Forgot your SecureCode?](#)

# Scheme

**Registration** User chooses a password of  $n$  characters from a set of  $N$

**Login** Challenge of  $m$  positions with response:

Positions:	1	2	3	4	5	6	7
User password:	a	s	h	u	f	1	0
Correct response:		s	h			1	

**Retry** In case of failure, user challenged again. Number of retries usually limited.

**Repeat** On next login, challenge changes.

# Motivations

Introduced for telephone banking: single observation by operator does not reveal whole secret.

Online, appears to impede several attacks:

- ▶ shoulder surfing
- ▶ key logging
- ▶ man-in-the-browser

Potentially, may also thwart:

- ▶ phishing
- ▶ offline attacks

Other attractions:

- ▶ easy extra authentication step (but not true 2FA)
- ▶ cheap (e.g., compared to hardware tokens)

# Origins

In UK banking: first introduced for telephone banking.

Matsumoto and Imai, *Human Identification Through Insecure Channel* (Eurocrypt '91). Related but more elaborate scheme:

- ▶ User has a password with known character set
- ▶ Challenge: word surrounded by detractor characters
- ▶ Response: *substituted* positions and detractors

Repeated several times.

Following work (e.g., Hopper & Bloom 2001): revised schemes and stronger guarantees, but showed required human computation steps are impractical.

So what about schemes actually in use?

# Questions

- ▶ What are the security assumptions behind current deployment of partial passwords?
- ▶ What are good choices for the system parameters: password length, character set size, challenge size?
- ▶ How many observations does an attacker need to learn whole password or answer next challenge?
- ▶ Are weak passwords such as dictionary words safe?
- ▶ Failure mode: should the challenge be changed after failed attempts?
- ▶ Are some challenge sequences better than others?
- ▶ How usable is the scheme?



# Survey

## Online banking survey: results

- ▶ Used widely in banks, online and telephone
- ▶ Elsewhere: credit cards, utilities, outside UK, . . .
- ▶ Usually part of a **multi-stage** authentication, alongside: names, user ids, account details, personal knowledge questions.
- ▶ Challenge sizes fixed, vary from 2-3 positions
- ▶ Challenge sequences appear random
- ▶ Mostly: ascending position challenges, no repeats
- ▶ Most repeat same challenge on retry
- ▶ Policies generally weaker than for full passwords

# Parameters

	character set size, $N$	password length, $n$	challenge size, $m$	second credential
Cooperative	10	4	2	question
ING DiBa (DE)	10	6	2	PIN
Tesco	10	6	2	password
Smile	10	6	2	question
Nationwide	10	6	3	password
AIB	10	5	3	question
B. of Ireland (IE)	10	6	3	date of birth
Nat West, step 1	10	4	2	pp, step 2
Nat West, step 2	36	6–20	3	pp, step 1
HBoS	36	6–15	3	password
3DSecure, Bol	36	8–15	3	credit card #
Standard Life	36	8–10	3	none
Skipton	36	8–30	3	question
First Direct	36	6–30	3	question
Barclays	52	6–8	2	PIN
HSBC (CA)	62	8	3	question

NB: snapshot from Sept. 2012. Thanks to Atif Hussain for help with survey.

# Guessing Attacks

## Mode of attack for guessing

- ▶ online attack against each account
- ▶ suppose a fixed number of attempts allowed:  $\beta$
- ▶ some background (e.g., dictionary), ideally limited
- ▶ no use of previous observations
- ▶ “trawling”: use best strategy on many accounts

Two typical instances of scheme:

*6 digit PIN*

- ▶  $N=10, n=6, m=2, \beta=6$

*8 character alphanumeric*

- ▶  $N=36, n=8, m=3, \beta=10$

## Guessing methods

1. **brute-force** (sample from uniform distribution)
2. **position-letter frequency** (ranked list per position)
3. **projection dictionary** (ranked list per challenge)
4. **dependent projection** (tree per challenge) *[later]*

Generate background tables by computation on:

- ▶ ordinary dictionary, e.g., /usr/share/dict/words
- ▶ dictionary with frequencies, e.g., RockYou

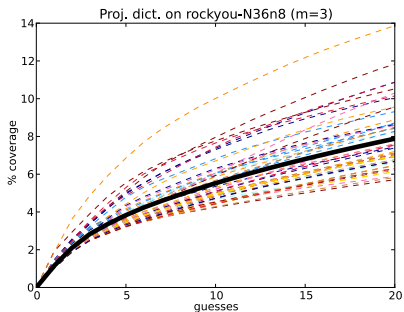
We calculate  **$\beta$ -success rate**: proportion of answers covered by the top  $\beta$  guesses.

# Example projection dictionary attack

<i>Challenge 2 3 6: Cum.%</i>	<i>Challenge 1 2 3: Cum.%</i>
1. a s o 1.10	1. i l o 1.29
2. l o y 1.98	2. p a s 2.42
3. r i e 2.79	3. m a r 3.40
4. 2 3 6 3.21	4. b a b 4.30
5. a r e 3.56	5. p r i 5.08

- ▶ The top 5 choices for two of the  $\binom{n}{m} = 56$  challenges
- ▶ Dictionary is RockYou (8-char alphanumeric) with frequencies
- ▶ 5.3m total, top 5 words in ranked dictionary covers 3.02%
- ▶ Top 5 full words:  
password, iloveyou, princess, 12345678, babygirl

# Example projection dictionary attack



- ▶ This shows the coverage of guesses for increasing  $\beta$
- ▶ Each line is a different challenge, bold is average
- ▶ Success rate for  $\beta=10$  is 5.5% versus 3.9% without projection



# Recording Attacks

## Mode of attack for recording

- ▶ online,  $\beta$  attempts per challenge, as before
- ▶ allow recording previous  $k$  challenge-response pairs

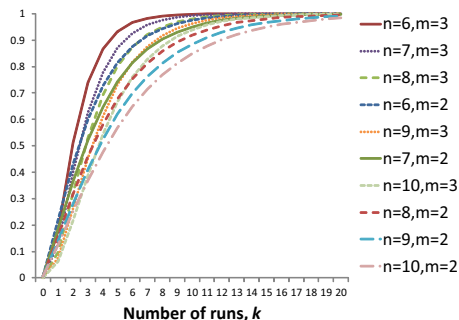
### Recording methods

1. **Pure recording**: only answer when positions known
2. **Recording+guessing**: guess remainder of positions

Combinatorics: we find equations for two different **success rates** for increasing  $k$ . They are probabilities of:

- ▶ answering the **next challenge**, or
- ▶ learning the **whole password**.

# Success rates for answering next challenge



This is a plot of

$$\sum_{j=0}^m \overline{s}_n^m(k, j) w_j$$

where  $0 \leq j \leq m$  positions are known in a challenge after  $k$  runs.

- ▶  $\overline{s}_n^m(k, j)$ : fraction of challenges with  $j$  known positions
- ▶  $w_j$ : the  $\beta$ -success rate for a particular guessing method

# Summary

## Results for typical parameters

Attack type	parameters	% success rate	
		PINs	alphanumeric
Brute force		6	0.002
Letter position	RockYou	17.2	0.3
Dictionary	RockYou	15.3	3.9
Proj. dictionary	RockYou	30.6	5.5
Recording	$k=1$ ( $k=4$ )	6.7 (63.1)	1.8 (59.0)
Recording + BF Guess	$k=1$ ( $k=4$ )	41.1 (83.8)	9.6 (69.1)
Recording + Best Dict	$k=1$ ( $k=4$ )	<b>60.2 (90.4)</b>	<b>25.2 (81.2)</b>

# Summary

- ▶ survey of partial password implementations
- ▶ model of partial password authentication scheme
- ▶ several attack methods, guessing and recording
- ▶ theoretical success rates measured analytically (pure recording) and empirically (using a dictionary)

## Future/ongoing work:

- ▶ Better attacks (dependent case)
- ▶ Unseen challenge (Goring et al, 2007)
- ▶ Failure modes, challenge schedule and format
- ▶ General study of multi-stage authentication
- ▶ Discuss more with banks. . .